

CS35X: Competitive Programming

Lecture 7: Stacks, Queues, Graphs

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Warmup Kattis Problem: relocation

**Problem debrief:
trianglesofasquare**

Stack, Queue ADTs

- Maintain ordered collection of items.
- Main Stack Operations:
 - **push(foo)** — add foo to top of stack
 - **pop()** — remove top element of stack
 - Items accessed in **LIFO** order
- Main Queue Operations:
 - **enqueue(foo)** — add foo to back of Queue
 - **dequeue()** — remove element from front of queue
 - Items accessed in **FIFO** order.
- Other Stack/Queue operation(s): ***isEmpty(), getSize()***
- Interfaces are ***simple***. Operations should be ***fast — O(1)!***

STL deque class

- **Double-ended queue** – add/remove items from front or back.
- `#include <deque>` // include pair library
- `deque<string> d;` // d is deque of strings
- `deque<int> q = {3,2,6};` // direct initialization
- `q.push_front(8);` // q is now [8,3,2,6]
- `q.push_back(9);` // q is now [8,3,2,6,9]
- `q.pop_front();` // q is now [3,2,6,9]
- `q.pop_back();` // q is now [3,2,6]
- `q.front();` // returns 3
- `q.back();` // returns 6
- `q.empty();` // returns false (q not empty)

Implementation Details

- Deque (**d**ouble-**e**nded **que**ues) built like *circular* ArrayLists.
- Adding/removing to front/end of deque is $O(1)$ in practice.
- **pop_front**, **pop_back** don't return the element.

Stack/Queue application: Graphs

- A graph $G = (V, E)$ is a set of *vertices* V along with a set of *edges* E .
- Graphs represent binary relationships.
 - Map: vertices == towns, edges == roads
 - Social network: vertices == users, edges == friendships
 - Temporal network: vertices == events, edge (u, v) : u happens before v .
- There are many interesting algorithms we can do on graphs

Graph Representation: Adjacency Lists

- Say graph has **n** vertices $0, \dots, n-1$.
- Represent graph **g** as array of *vector<int>*:
 - **g[i]**: vector of neighbors of i .

A first Graph algorithm: BFS

- Visit all nodes in graph, using queue to manage exploration
- `deque<int> q;`
- `q.push_back(s);`
- `visited[s] = true;`
- `while(!q.empty()) {`
 - `v = q[0];`
 - `q.pop_front();`
 - `for(int i=0; i<g[v].size(); i++) {`
 - `u = g[v][i];`
 - `if(! visited[u]) {`
 - `// visit u`
 - `q.push_back(u);`
 - `}`
 - `}`
- `}`

Things we can do with BFS:

- Are all vertices connected?
- Find shortest length path from $s \rightarrow t$.
- Identify vertices reachable from s .
- Count # connected *components* of graph
- ...

Graph Implementation tips

- Represent vertices of graph as integers.
- Adjacency List can be dictionary instead of array, but **huge** time penalty.
- Adjacency List can store:
 - Vertex neighbors
 - Edges out of a vertex
- Create a helper function that loads graph into an adj list
- Consider defining a class to represent *weighted* edges.
- String vertices:
 - Store Dictionary mapping integers to (string) vertex label
 - Run graph algorithm on graph using int vertices
 - Use dictionary to recover vertex names if needed

**Kattis Problem:
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