

7.1 binary trees and searching

Tuesday, October 18, 2022

Today:

- 20 questions game
- introduction of trees, binary trees, and binary search trees
- (if time) dictionary ADT, keys, values

Reminder: lab 6 is due Oct 26th (next week) and has three parts:

1. implement stacks and queues
2. solve mazes using BFS and DFS
3. induction proofs

You must come to lab this week unless your team has completed and pushed all parts.

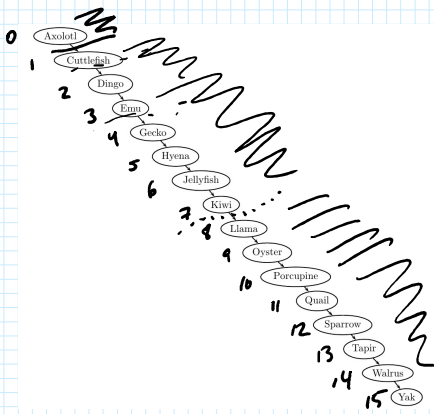
Think of your favorite animal.

I have a list of animals I can think of:

- axolotl
- yak
- cuttlefish
- walrus
- dingo
- tapir
- emu
- sparrow
- gecko
- quail
- hyena
- porcupine
- jellyfish
- oyster
- kiwi
- llama

(I'm storing this in an array list or LinkedList, probably.)

How can I get better (faster) at finding if your favorite animal is on my list? (by changing my list or my strategy)



Binary search for X

- ① If X is in the list, it must be somewhere 0...15
 $\lfloor \frac{0+15}{2} \rfloor = 7$ array[7] = kiwi
 X is alphabetically < kiwi
- ② If X is in the list, it's in 0...6
 $\frac{0+6}{2} = 3$ array[3] = Emu
 X is < Emu
- ③ If X is in the list, it's in 0...2
 $\frac{0+2}{2} = 1$ array[1] = Cuttlefish
 X is < Cuttlefish *
- ④ If X is in the list, it's in 0...1

⑤ If X is in the list, it's after index 0 before index 1 return "X not found!"

Once the data is in sorted order, I can use $\lfloor \frac{0+1}{2} \rfloor = 0$ array[0] = axolotl X > axolotl *

binary search:

idea: repeatedly reduce the size of search space by half (until you find element or determine it's not there)

Worst case takes $O(\log_2(n))$ where n = size of list

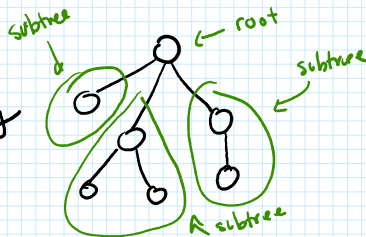
$$\log_2(1 \text{ million}) \approx 20$$

$$\log_2(1 \text{ billion}) \approx 30$$

So far all our data structures have been LINEAR.
 We'll use the idea of binary search to build
 our first HIERARCHICAL data structure: a TREE.

TREE:

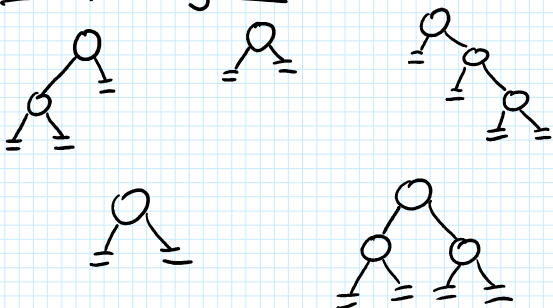
- Collection of nodes
- in general, each node may have any number of branches
- recursive structure



We will focus on BINARY TREES:

either: empty
 or: consists of a node that contains links to 2 trees below it, called left and right

example binary trees:



(Usually we won't bother to draw all the null pointers.)

Terminology

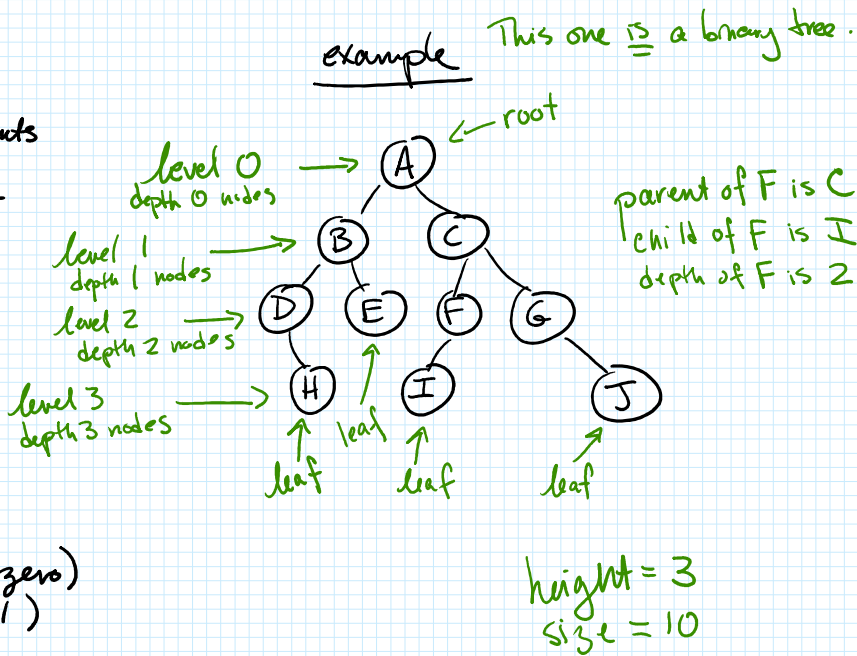
node

contains { data ← called "key"
 { pointers to left & right subtrees

children (of a node) the nodes directly linked below

(pointers to left & right subtrees)

- children (of a node) the nodes directly linked below
- parent (of a node) the unique node directly above
- root (of a tree) the unique node with no parents
- size (of a tree) number of nodes in the tree
- leaf a node with no children
- depth (of a node) how far node is from root (number of links you have to traverse) (root is depth zero)
- height (of a tree) the max depth of any node in the tree (tree with only root \rightarrow height zero) (empty tree \rightarrow height -1)



If we just put all animals in the tree structure arbitrarily, that won't help make the search more efficient.

We need a special kind of tree:

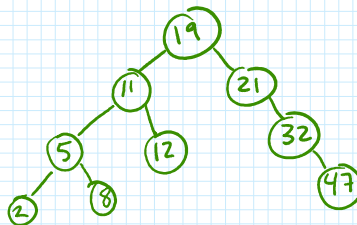
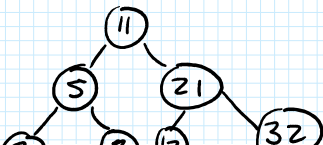
BINARY SEARCH TREES!

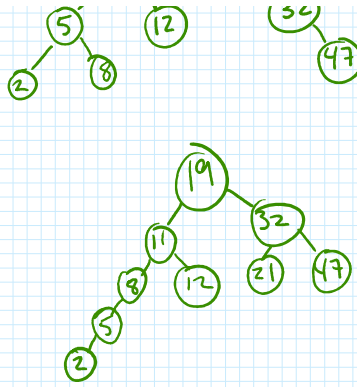
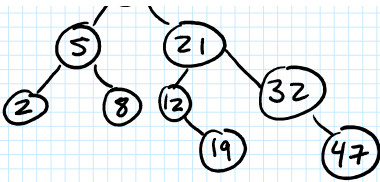
- a binary tree has a special property governing all keys stored within nodes, property is true at all nodes & helps with searching

BINARY SEARCH TREE PROPERTY

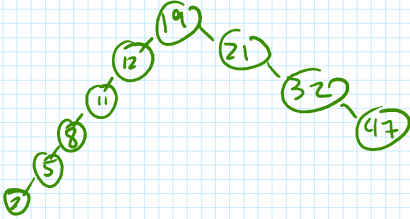
- All keys in the left subtree of a node must be less than the key at that node.
- All keys in the right subtree of a node must be greater than the key at that node.

example BST:

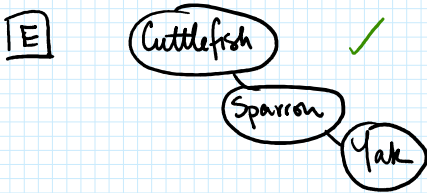
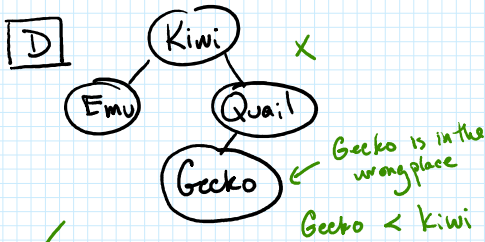
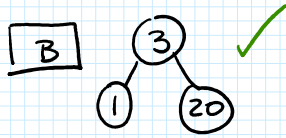
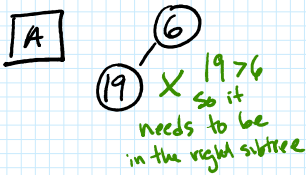




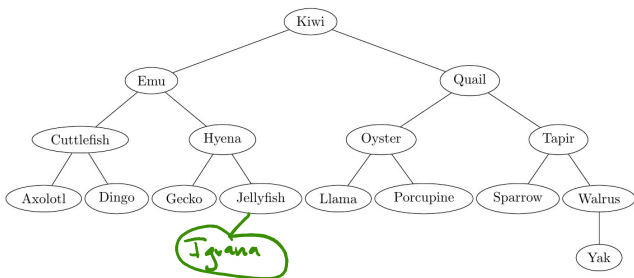
Make another BST with 19 at the root:
2, 5, 8, 11, 12, 19, 21, 32, 47



Which of these are valid BSTs?



Let's do our binary search again,
but this time with a BST:



Kiwi? no, after
Quail? no, before
Oyster? no, after
Porcupine? yes!

Note: search takes $O(\text{height of tree})$ many questions.

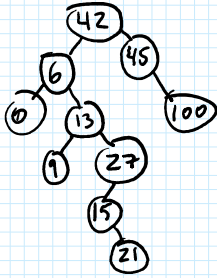
How to add a new key to the tree?

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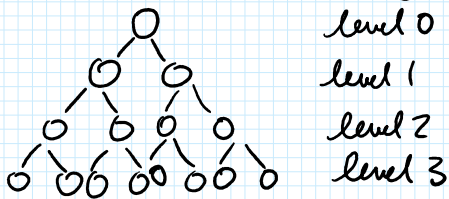
ex: Iguana: do the same steps as search, add a new node once we reach an empty tree.

Create a BST from these keys by inserting

them in order: 42, 6, 13, 45, 27, 9, 15, 9, 21, 100



Best possible structure is a fully-packed tree:



nodes = 15

$$2^{(\text{height} + 1)} - 1 = \# \text{ nodes in a full binary tree of that height}$$

If the tree has n nodes:

best case height: $O(\log_2(n))$ fully packed

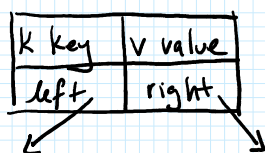
worst case height: $O(n)$ is basically a linked list

Questions to think about:

- If we don't know all the keys beforehand, how do we keep the BST balanced as we insert new keys?
- What should we do if we want to remove a key from the BST?

implementation details:

Linked BST Node



} contains 4
data
members

```
private: // data
    K key
    v value
    Linked BST Node < k, v > * left
    Linked BST Node < k, v > * right
public: // methods
    get Key, set Key
    get Value, set Value
    get Left, set Left
    get Right, set Right
```

Linked BST

private: //data

Linked BSTNode<k,v> * root

int size

Note: you can reach the entire tree
by starting at the root.